# Characterizing homomorphism functions and the limit theory of graphs

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Mike Freedman: Which graph parameters are partition functions?

Jennifer Chayes: What is the limit of a randomly growing graph sequence like the internet?

Vera Sós: How to characterize generalized quasirandom graphs?

Homomorphism: adjacency-preserving map hom(F,G) := # of homomorphisms of F into G $t(F,G) = \frac{hom(F,G)}{|V(G)|^{|V(F)|}} \leftarrow Probability \text{ that random map}} V(F) |V(C)| = h = h$ 

Weighted version:

$$H = (V, E, \alpha, \beta), \quad \alpha : V \to \mathbf{i}_{+}, \quad \beta : E \to \mathbf{i}$$
$$\hom(G, H) := \sum_{\varphi : V(G) \to V(H)} \prod_{i \in V(G)} \alpha_{\varphi(i)} \prod_{i j \in E(G)} \beta_{\varphi(i)\varphi(j)}$$

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### k-labeled graphs

*k*-labeled graph: *k* nodes labeled 1,...,*k*, any number of unlabeled nodes



 $G_1, G_2$ : *k*-labeled graphs  $G_1G_2 := G_1 \stackrel{\bullet}{\mathbf{E}} G_2$ , labeled nodes identified

#### Graph parameter

# Graph parameter: isomorphism-invariant function *f* on finite graphs





#### Which parameters are homomorphism functions?



Generalizations, analogues and relaxations

Generalizations: directed graphs (Schrijver), hypergraphs, semigroups (L-Schrijver)

Analogues: edge coloring models (Szegedy), unweighted graphs (L-Schrijver), tensor algebras (Schrijver), categories

Which graph parameters have connection matrices - positive semidefinite?

- finite rank?

#### Locally finite categories

 $\mathcal{K}$ : category,  $\mathcal{K}(a,b)$  finite for any two objects

(other conditions: zero object, generator object, pushouts, epi-mono decompositions)



# Locally finite categories

f: objects 🐼 reals

M(f,a): matrix indexed by morphisms from object a $M(f,a)_{\alpha\beta} = f(\alpha \wedge \beta)$ 



#### Corollary for simple graphs



 $\mathbf{W}_{0} = \{ W : [0,1]^{2} \otimes [0,1] \text{ symmetric, measurable} \}$ 

(graphons)



Graphs [X] Graphons





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 $A_{G}$ 

 $\mathbf{W}_{0} = \{ W : [0,1]^{2} \otimes [0,1] \text{ symmetric, measurable} \}$ (graphons)

$$t(F,W) = \bigcup_{[0,1]^{V(F)}} \widetilde{\mathbf{O}}_{E(F)} W(x_i, x_j) dx$$

$$t(F, W_G) = t(F, G)$$

 $(G_1, G_2, ...)$  convergent: " $F t(F, G_n)$  is convergent Borgs, Chayes, L, Sós, Vesztergombi

# $G_n \otimes W$ : "F $t(F,G_n) \otimes t(F,W)$

This is for dense graphs
(M) mean field theory

For every convergent graph sequence  $(G_n)$  there is a graphon W such that  $G_n \boxtimes W$ .

L-Szegedy

For every graphon W there is a graph sequence  $(G_n)$  such that  $G_n \boxtimes W$ . L-Szegedy

*W* is essentially unique (up to measure-preserving transformation).

**Borgs-Chayes-Lovasz** 

# Semidefinite connection matrices

f: graph parameter



L - Szegedy

#### Finite rank connection matrices

If M(f,k) has finite rank, then f(G) can be evaluated in polynomial time for graphs with tree-width at most k.

L-Welsh

M(f,k) has finite rank for many parameters: hom(.,H), hom(F,.), # of Hamilton cycles, max no. of independent nodes,  $1_{planar}$ ,  $1_{P}$  for any minor-closed property P.

If f can be defined by a monadic second order counting problem, then M(f,k) has finite rank.

Godlin-Kotek-Makowski



#### *k*-labeled quantum graph:

$$x = \mathop{\mathbf{a}}_{G} x_{G} \mathcal{A}_{G} \mathcal{A}_{G}$$
 finite formal sum of *k*-labeled graphs



Q<sub>k</sub> = {k-labeled quantum graphs}

$$\begin{array}{c} \overbrace{G}^{\bullet} & x_G G \overbrace{G}^{\bullet} & y_G G \overbrace{G}^{\bullet} & = \\ \overbrace{G}^{\bullet} & a_{G_1,G_2} & x_{G_1} y_{G_2} G_1 G_2 \end{array} \\ Q_k \text{ is a commutative algebra with unit element} \end{array}$$

Inner product:  $\langle x, y \rangle = f(xy)$ 

$$N_{k}(f) = \{x \boxtimes Q_{k} : \langle x, y \rangle = 0 \boxtimes y\}$$

 $Q_k/N_k(f)$  is a commutative algebra

 $Q_k/N_k(f)$  is finite dimensional iff M(f,k) has finite rank

<...> is semidefinite iff M(f,k) is positive semidefinite

 $p_1, \ldots, p_q$ : idempotent basis in  $Q_k$  nodes (states)

$$[\mathbf{M}] = p_1 + \dots + p_n$$
  
$$< [\mathbf{M}], p_i > \text{ nodeweights}$$
  
$$< [\mathbf{M}] - [\mathbf{M}], p_i [\mathbf{M}] p_j = \text{signweights}$$

